

#### **Outline**

- Why FOL?
- Syntax and semantics of FOL
- Using FOL
- Wumpus world in FOL
- Knowledge engineering in FOL

# Pros and cons of propositional logic

- Propositional logic is declarative
- Propositional logic allows partial/disjunctive/negated information
- Propositional logic is compositional:
  - meaning of  $B_{1,1} \wedge P_{1,2}$  is derived from meaning of  $B_{1,1}$  and of  $P_{1,2}$
- Meaning in propositional logic is context-independent
  - (unlike natural language, where meaning depends on context)
- ® Propositional logic has very limited expressive power
  - (unlike natural language)
  - E.g., cannot say "pits cause breezes in adjacent squares"
    - except by writing one sentence for each square

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# First-order logic

- Whereas propositional logic assumes the world contains facts,
- first-order logic (like natural language) assumes the world contains
  - Objects: people, houses, numbers, colors, baseball games, wars, ...
  - Relations: red, round, prime, brother of, bigger than, part of, comes between, ...
  - Functions: father of, best friend, one more than, plus, ...

## Syntax of FOL: Basic elements

- Constants KingJohn, 2, NUS,...
- Predicates Brother, >,...
- Functions Sqrt, LeftLegOf,...
- Variables x, y, a, b,...
- Connectives  $\neg$ ,  $\Rightarrow$ ,  $\wedge$ ,  $\vee$ ,  $\Leftrightarrow$
- Equality =
- Quantifiers ∀,∃

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#### **Atomic sentences**

Atomic sentence =  $predicate (term_1,...,term_n)$ 

or  $term_1 = term_2$ 

Term =  $function (term_1,...,term_n)$ 

or constant or variable

E.g.,

- Brother(KingJohn,RichardTheLionheart)
- > (Length(LeftLegOf(Richard)), Length(LeftLegOf(KingJohn)))

## **Complex sentences**

 Complex sentences are made from atomic sentences using connectives

$$\neg S$$
,  $S_1 \land S_2$ ,  $S_1 \lor S_2$ ,  $S_1 \Rightarrow S_2$ ,  $S_1 \Leftrightarrow S_2$ ,

E.g.

Sibling(KingJohn,Richard) ⇒
 Sibling(Richard,KingJohn)

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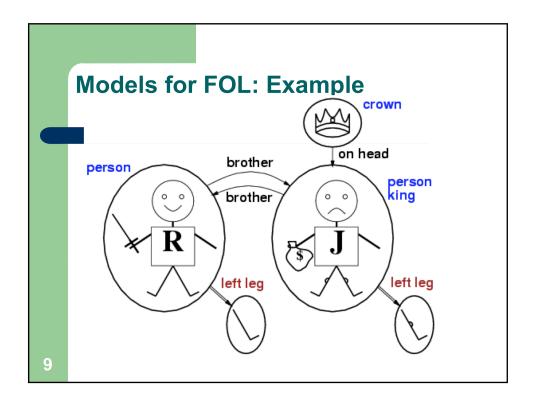
# **Truth in first-order logic**

- Sentences are true with respect to a model and an interpretation
- Model contains objects (domain elements) and relations among them
- Interpretation specifies referents for

 $\begin{array}{ccc} \text{constant symbols} & \to & \text{objects} \\ \text{predicate symbols} & \to & \text{relations} \end{array}$ 

function symbols → functional relations

An atomic sentence predicate(term<sub>1</sub>,...,term<sub>n</sub>) is true iff the objects referred to by term<sub>1</sub>,...,term<sub>n</sub> are in the relation referred to by predicate



## **Universal quantification**

- ∀<variables> <sentence>
- Everyone at SMU is smart:

 $\forall x \ At(x,SMU) \Rightarrow Smart(x)$ 

- $\forall x P$  is true in a model m iff P is true with x being each possible object in the model
- Roughly speaking, equivalent to the conjunction of instantiations of P

```
At(KingJohn,SMU) ⇒ Smart(KingJohn)
```

 $\wedge$  At(Richard,SMU)  $\Rightarrow$  Smart(Richard)

 $\wedge$  At(Rupert,SMU)  $\Rightarrow$  Smart(Rupert)

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#### A common mistake to avoid

- Typically, ⇒ is the main connective with ∀
- Common mistake: using ∧ as the main connective with ∀:

 $\forall$  x At(x,SMU)  $\land$  Smart(x) means "Everyone is at SMU and everyone is smart"

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## **Existential quantification**

- ∃<variables> <sentence>
- Someone at SMU is smart:
- ∃x At(x,SMU) ∧ Smart(x)
- $\exists x P$  is true in a model m iff P is true with x being some possible object in the model
- Roughly speaking, equivalent to the disjunction of instantiations of P
  - At(KingJohn,SMU) ∧ Smart(KingJohn)
  - ∨ At(Richard,SMU) ∧ Smart(Richard)
  - ∨ At(Rupert,SMU) ∧ Smart(Rupert)

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#### Another common mistake to avoid

- Typically, ∧ is the main connective with ∃
- Common mistake: using ⇒ as the main connective with ∃:

 $\exists x \, \mathsf{At}(\mathsf{x}, \mathsf{SMU}) \Rightarrow \mathsf{Smart}(\mathsf{x})$ 

is true if there is anyone who is not at SMU!

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## **Properties of quantifiers**

- ∀x ∀y is the same as ∀y ∀x
- ∃x∃y is the same as∃y∃x
- $\exists x \forall y \text{ is not the same as } \forall y \exists x$
- ∃x ∀y Loves(x,y)
  - "There is a person who loves everyone in the world"
- ∀y∃x Loves(x,y)
  - "Everyone in the world is loved by at least one person"
- Quantifier duality: each can be expressed using the other
- $\forall x \text{ Likes}(x, \text{IceCream}) \quad \neg(\exists x \neg \text{Likes}(x, \text{IceCream}))$
- $\exists x \text{ Likes}(x, \text{Broccoli})$   $\neg(\forall x \neg \text{Likes}(x, \text{Broccoli}))$

## **Equality**

- term<sub>1</sub> = term<sub>2</sub> is true under a given interpretation if and only if term<sub>1</sub> and term<sub>2</sub> refer to the same object
- E.g., definition of *Sibling* in terms of *Parent*:

```
\forall x,y \; Sibling(x,y) \Leftrightarrow [\neg(x = y) \land (\exists m,f \neg (m = f) \land Parent(m,x) \land Parent(f,x) \land Parent(m,y) \land Parent(f,y))]
```

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## **Using FOL**

The kinship domain:

- Brothers are siblings
   ∀x,y Brother(x,y) ⇔ Sibling(x,y)
- One's mother is one's female parent

   ∀m,c Mother(c) = m ⇔ (Female(m) ∧ Parent(m,c))
- "Sibling" is symmetric
   ∀x,y Sibling(x,y) ⇔ Sibling(y,x)

#### **Interacting with FOL KBs**

 Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at t=5.

Tell(KB,Percept([Smell,Breeze,None],5))Ask(KB, $\exists$ a BestAction(a,5))

- i.e., does the KB entail some best action at *t*=5?
- Answer: Yes, {a/Shoot} ← substitution (binding list)
- Given a sentence S and a substitution σ,
- $S\sigma$  denotes the result of plugging  $\sigma$  into S; e.g.,

S = Smarter(x,y)

 $\sigma = \{x/Hillary, y/Bill\}$ 

 $S\sigma = Smarter(Hillary,Bill)$ 

• Ask(KB,S) returns some/all  $\sigma$  such that  $KB \models \sigma$ 

# Knowledge base for the wumpus world

- Perception
  - ${}- \ \forall \, t, s, b \, \, \mathsf{Percept}([s, b, \mathsf{Glitter}], t) \Rightarrow \mathsf{Glitter}(t)$
- Reflex
  - ∀t Glitter(t) ⇒ BestAction(Grab,t)

#### **Deducing hidden properties**

∀x,y,a,b Adjacent([x,y],[a,b]) ⇔
 [a,b] ∈ {[x+1,y], [x-1,y],[x,y+1],[x,y-1]}

Properties of squares:

•  $\forall$  s,t At(Agent,s,t)  $\land$  Breeze(t)  $\Rightarrow$  Breezy(s)

#### Squares are breezy near a pit:

- Diagnostic rule---infer cause from effect
   ∀s Breezy(s) ⇒ \Exi{r} Adjacent(r,s) ∧ Pit(r)\$
- Causal rule---infer effect from cause
   ∀r Pit(r) ⇒ [∀s Adjacent(r,s) ⇒ Breezy(s)\$]

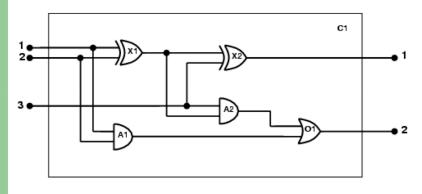
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# Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and constants
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

#### The electronic circuits domain

#### One-bit full adder



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## The electronic circuits domain

- 1. Identify the task
  - Does the circuit actually add properly? (circuit verification)
- 2. Assemble the relevant knowledge
  - Composed of wires and gates; Types of gates (AND, OR, XOR, NOT)
  - Irrelevant: size, shape, color, cost of gates
- 3. Decide on a vocabulary
  - Alternatives: Type(X<sub>1</sub>) = XOR Type(X<sub>1</sub>, XOR) XOR(X<sub>1</sub>)

#### The electronic circuits domain

- 4. Encode general knowledge of the domain
  - $\forall$  t<sub>1</sub>,t<sub>2</sub> Connected(t<sub>1</sub>, t<sub>2</sub>) ⇒ Signal(t<sub>1</sub>) = Signal(t<sub>2</sub>)
  - $\forall$ t Signal(t) = 1  $\lor$  Signal(t) = 0
  - \_ 1 ≠ 0
  - $\forall t_1, t_2$  Connected( $t_1, t_2$ )  $\Rightarrow$  Connected( $t_2, t_1$ )
  - ∀g Type(g) = OR ⇒ Signal(Out(1,g)) = 1 ⇔ ∃n
     Signal(In(n,g)) = 1
  - $\forall$  g Type(g) = AND  $\Rightarrow$  Signal(Out(1,g)) = 0  $\Leftrightarrow$  ∃n Signal(In(n,g)) = 0
  - ∀g Type(g) = XOR ⇒ Signal(Out(1,g)) = 1 ⇔ Signal(In(1,g))
     ≠ Signal(In(2,g))
  - $\forall$  g Type(g) = NOT ⇒ Signal(Out(1,g)) ≠ Signal(In(1,g))

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#### The electronic circuits domain

5. Encode the specific problem instance

```
Type(X_1) = XOR Type(X_2) = XOR

Type(A_1) = AND Type(A_2) = AND
```

Type( $O_1$ ) = OR

 $\begin{aligned} & \text{Connected}(\text{Out}(1, X_1), \text{ln}(1, X_2)) & \text{Connected}(\text{ln}(1, C_1), \text{ln}(1, X_1)) \\ & \text{Connected}(\text{Out}(1, X_1), \text{ln}(2, A_2)) & \text{Connected}(\text{ln}(1, C_1), \text{ln}(1, A_1)) \end{aligned}$ 

Connected(Out(1, $A_2$ ),In(1, $O_1$ )) Connected(In(2, $C_1$ ),In(2, $X_1$ )) Connected(Out(1, $A_1$ ),In(2, $O_1$ )) Connected(In(2, $C_1$ ),In(2, $A_1$ ))

Connected(Out(1, $X_2$ ),Out(1, $C_1$ )) Connected(In(3, $C_1$ ),In(2, $X_2$ ))

Connected(Out(1,O<sub>1</sub>),Out(2,C<sub>1</sub>)) Connected(In(3,C<sub>1</sub>),In(1,A<sub>2</sub>))

#### The electronic circuits domain

6. Pose queries to the inference procedure: What are the possible sets of values of all the terminals for the adder circuit?

```
\exists i_1, i_2, i_3, o_1, o_2 \text{ Signal}(\text{In}(1, C_1)) = i_1 \land \text{Signal}(\text{In}(2, C_1)) = i_2 \land \text{Signal}(\text{In}(3, C_1)) = i_3 \land \text{Signal}(\text{Out}(1, C_1)) = o_1 \land \text{Signal}(\text{Out}(2, C_1)) = o_2
```

7. Debug the knowledge baseMay have omitted assertions like 1 ≠ 0

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#### **Summary**

- First-order logic:
  - objects and relations are semantic primitives
  - syntax: constants, functions, predicates, equality, quantifiers
- Increased expressive power: sufficient to define wumpus world