

Brief Announcement: Opportunistic Information Dissemination in Mobile Ad-hoc Networks:

adaptiveness vs. obliviousness and randomization vs. determinism

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Mobile Ad-hoc Network (MANET)

- Mobile set of nodes (processors with radio)
- No stable communication infrastructure
- Multihop network

E.g.



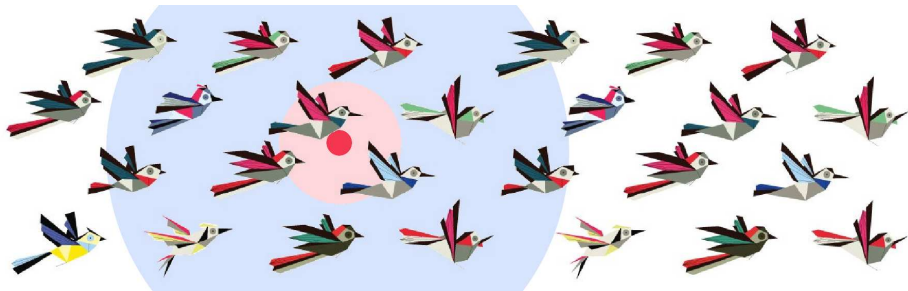
Opportunistic Communication

Thanks to mobility and asynch activation
communication between x and y is feasible
even if a path never exists! (a *chrono-path*)

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The Dissemination Problem

Some information held by a given source node x at time t ,
must be disseminated to some set of nodes $S \subset V$.



In order to prove lower bounds we use **Geocast**.

Model

- Network:
 - n mobile nodes deployed in \mathbb{R}^2
 - slotted time steps:
 - slot length dominated by communication time
 - same for all nodes
- Node:
 - unique ID in $[n]$
 - may start/fail at any time slot
 - radio communication:
 - unique radio channel \implies collisions
 - background noise \equiv collision noise \implies no collision detection
 - no simultaneous reception & transmission
 - limited range $r \implies$ multihop network

Model

- Adversary:

- initial position and movement
- de/activation schedule (lower bounds don't use it!)

limited by three parameters:

- a maximum speed $v_{\max} > 0$
- the system must be (α, β) -**connected**, $\alpha, \beta \in \mathbb{Z}^+$

Definition ((α, β)-connectivity)

While moving at $\leq v_{\max}$ speed, \forall non-trivial partition (S, \overline{S}) ,

$\exists \leq \alpha$ consecutive steps without a β -**stable edge** between S and \overline{S} .

(an edge is k -stable at time t if it exists for k consecutive steps $[t, t + k - 1]$)

Model

(α, β) -connectivity, for the partition defined by the information

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Model

- Randomized Protocols:
 - *oblivious* [C'01]: protocol access sequence of random variables at each node, independent of execution and mutually independent.
 - *locally adaptive*: same but rv's may be mutually dependent. (still independent of the execution)
 - *fair* [C'01]: all nodes transmit with same probability in any given time step. (orthogonal def)

Results

		randomized	deterministic [FMMZ'10]
l.b.	oblivious	$w.p. \geq 2^{-n/2} \Rightarrow \Omega(\alpha n + n^2 / \log n)$	$\Omega(\alpha n + n^3 / \log n)$
	adaptive	$\Omega(\alpha n + n^2 / \log n)$ exp.	$\Omega(\alpha n + n^2)$
	fair	$w.p. \geq 2^{-n/2} \Rightarrow \Omega(\alpha n + n^2 / \log n)$	–
u.b.	oblivious	$O(\alpha n + (1 + \alpha/\beta) n^2 / \log n)$ w.h.p.	$O(\alpha n + n^3 \log n)$
	adaptive	–	$O(\alpha n + n^2)$
	fair	$O(\alpha n + (1 + \alpha/\beta) n^2 / \log n)$ w.h.p.	–

Conclusions

- local adaptiveness $\Omega\left(\alpha n + \frac{n^2}{\log n}\right) \exp$
does not help w.r.t. obliviousness $O\left(\alpha n + \left(1 + \frac{\alpha}{\beta}\right) \frac{n^2}{\log n}\right) \text{ whp.}$
- linear separation between oblivious randomized $O\left(\alpha n + \left(1 + \frac{\alpha}{\beta}\right) \frac{n^2}{\log n}\right) \text{ whp}$
and oblivious deterministic $\Omega\left(\alpha n + \frac{n^3}{\log n}\right).$
- log separation between adaptive randomized $O\left(\alpha n + \left(1 + \frac{\alpha}{\beta}\right) \frac{n^2}{\log n}\right) \text{ whp}$
and adaptive deterministic $\Omega\left(\alpha n + n^2\right).$

Thank you