



THE START

Quantum Key Distribution [QKD]

Quantum Keys By Polarization (w/) or (w/o) Entanglement

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$$\frac{d^2 ()}{dx^2}$$

2. MOTIVATION

All encryption methods fail in one way. The distribution of keys is **insecure**.

All current encryption methods can be broken with enough computing power [which is increasing] except one-time-pads **OTP** (see above).

QKD distribution is secure.
QKD can't be broken (it is a OTP).

3. Encryption – Why Bother With QKD? & Current State of QKD & Overview of QKD

Why Bother With QKD? (Current Problems Intro)

- A Current State of Network Encryption Usage**
- B Why should we understand it?**
- C Potential Weaknesses**
- D What is The Classical Key Distribution Problem?**
- E Therefore QKD**

A. Current State of Network Encryption Usage

- **DES → AES (Symmetric all)**
- **PKI (Asymmetric – 2 key)**
- **IPSEC (Symmetric, AES)**
- **SSL/TLS (Negotiated – PKI/Symmetric)**

B. Why should we understand it?

- **PKI & AES**

- **Founded on a conjecture**
 - **Ultimately on computational burden too large for an enemy.**
- **BUT Faster Computers Coming**
- **BUT Quantum Computers**
Will Break It!

C. Potential Weaknesses

1. **PKI Based on assumptions about factoring large numbers.**
2. **Security of 3rd Party PKI Key Vault**
3. **Security of (P_r) Key**
 - *Distribution Problem*
 - *3rd Party Vault Current MO*

D. What is The Classical Key Distribution Problem?

- **One-time-pads (1 key/1 message)**
IT IS SYMMETRIC! [You & Me: Pads \equiv]
Key as long as message ($|K| \sim |M|$)
Vernam Cipher ($K \oplus M$) [XOR]
ONLY Unbreakable Code
(Proved by Shannon)
If done right.
The Key (PAD)
Distribution Problem Again

E. WHY QKD? [1/3]

- **Computation capability is increasing non-linearly**
- **Quantum Computers Promise to Completely Negate Efficacy of Current Encryption Technology (i.e., *kill it dead*) (not imminent)**

E. WHY QKD? [2/3]

- **QKD is Based On Physics**
- **Unaffected By Either:**
 - **Current Computer Technology**
 - or**
 - **QUANTUM COMPUTING CAPABILITY**
- **It is a handshake protocol**
- **It can **sense Eve** (*Alice, Eve, Bob*)**
- **After Key Distribution:**
 - **Use classical**
 - or**
 - **Q-encryption**

E. WHY QKD? [3/3]

**QKD SOLVES THE
KEY DISTRIBUTION
PROBLEM
&
IS UNBREAKABLE**

Summary: So What?

- **More Secure Data Transmission**
- **QKD Used For:**
 - **IPSEC (for Internet) (& SSL)**
 - **Replace PKI, AES**
 - **It is a Vernam One-Time-Pad (Unbreakable!)**
 - **Solves the key distribution problem**
 - **Borming's Dissertation (for Grids)**
 - **Chinese from a satellite.**
 - **Chinese national effort to secure networking.**

Current QKD State

- **QKD**
 - There are products that do it (100+ km distances) [MAGIQTECH]
 - Open air QE coming to a satellite near you
 - BBN Boston Network & Vienna Network
- **QKD In TCP/IP**
 - Research progressing
- **QKD Education**
 - QE appearing in CS texts [Tanenbaum's Networking]
- **Cultural Motivation to Learn**
 - 30% GDP derives from QM
[Waite, Stephen R., 2002]

OVERVIEW of QKD [1/2]

A Crypto Key

- A key for encryption/decryption is sent using Quantum Mechanical Phenomena.
- The key may be a quantum encryption key or a non-Quantum encryption Key, e. g., a PKI private key.
- The transmission may or may not involve entanglement.

Entanglement

Two (or more) particles created as single coupled complimentary set. A measurement of one determines the complimentary value of the other(s) regardless of separation.

OVERVIEW of QKD [2/2]

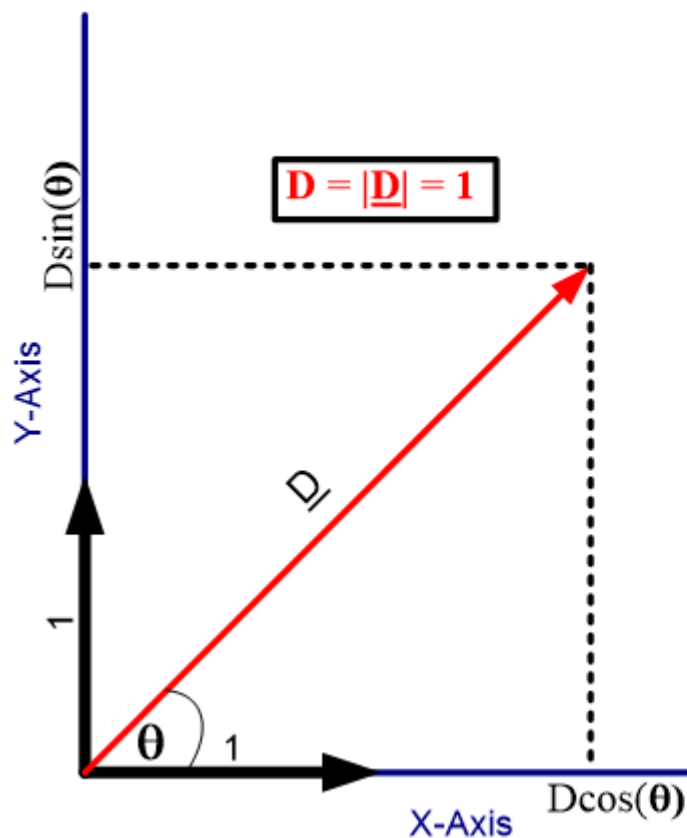
One-time-pads (1TP)

- **QKD is to used create a shared key for a 1TP**
- **The 1TP is used to send an encrypted message**
- **Only proved unbreakable encryption scheme.**
- **This is done many time/sec (>100)**

4. Vectors & QM

Vectors

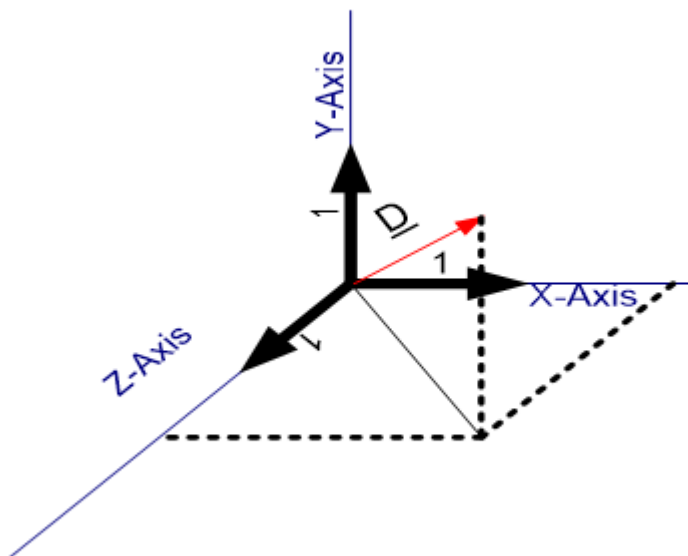
TRICK



Notice: *2-D Special Case*

If we know the X component of vector \underline{D} of length = 1, is $\underline{0}$ along the X axis, we know \underline{D} lies in the Y axis. **UNIQUENESS** (up to \pm). Not so in 3-D or greater.

Vectors



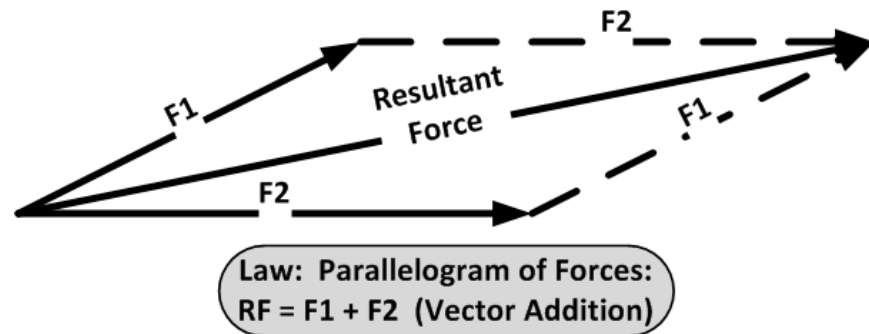
Notice:

If we know the vector $|\underline{D}| = 1$ has component of $\underline{0}$ length along the X axis, **ALL** we know it lies in the Y-Z plane. There are ∞ vectors (or lines) in Y-Z orthogonal to X. **NON-UNIQUENESS.**

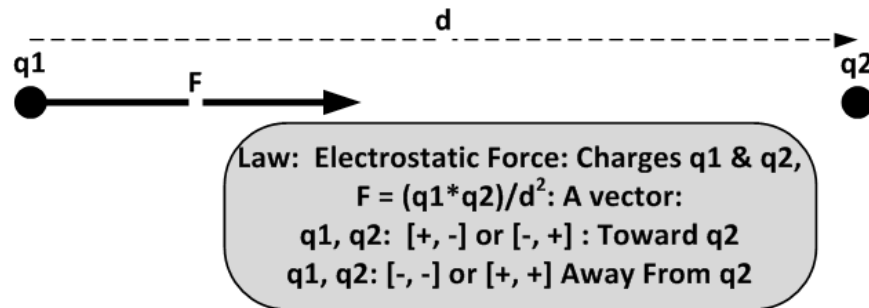
Vectors

Vector Examples

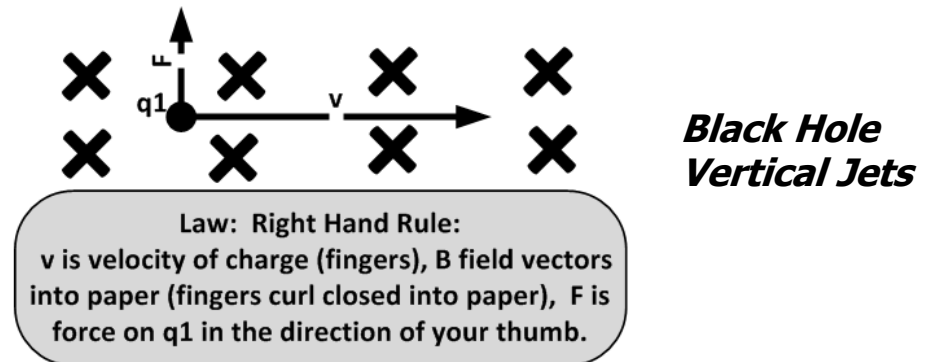
Mechanical FORCE



Electric FORCE



Magnetic FORCE



WHY Vectors (Linear Algebra)?

The **world** can be **effectively** modelled by Objects that have Observable States with Measureable [real] Values with known probabilities of measurement.

Objects can be **effectively** modelled by Hermitian Operators on vectors.

States can be **effectively** modelled by vectors [combination of eigenvectors].

Measureable values of the object can be **effectively** modelled by eigenvalues of the eigenvectors of the object. Hermitian \Rightarrow real.

WHY Vectors (Linear Algebra)?

The **probability** of finding the initial [before measurement] system in the final eigenstate vector [i] with measured eigenvalue [λ] after measurement, can be **effectively** modelled as the square of the projection of the initial [before measurement] system vector onto the eigenvector [i] found as the result of the measurement.

Example (Continued):

$$H = \lambda \underline{v}$$

$$H = \sum_{i=1}^{i=N} \lambda_i (\underline{v}_i \underline{v}_i^T) == \sum_{i=1}^{i=N} \lambda_i P[\underline{v}_i]$$

Hermitian operators have a spanning set of eigenvectors with all eigenvalues real.

https://en.wikipedia.org/wiki/Spectral_theorem

For convenience we take the eigenvectors & other state vectors to be of *unit length*.

If the operator is real symmetric this is the principle components theorem.

https://en.wikipedia.org/wiki/Principal_component_analysis

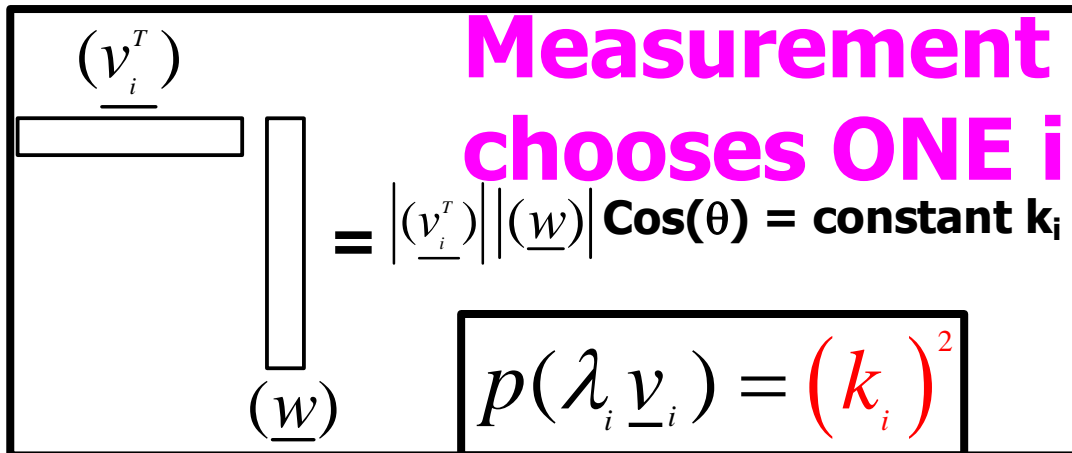
$$[A = X^T X] \Rightarrow A^T = (X^T X)^T = X^T X = A$$

Example (Continued):

Projector on the eigenvector.

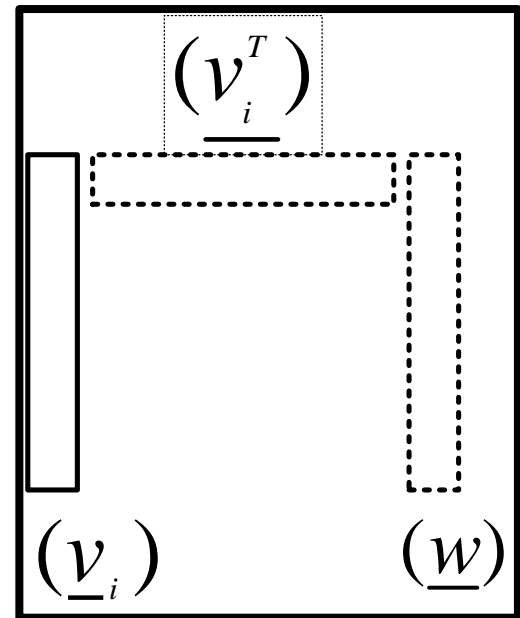
$$(\underline{v}_i \underline{v}_i^T) \underline{w} = \underline{v}_i (\underline{v}_i^T \underline{w}) = P[\underline{v}_i] \underline{w} = \underline{v}_i k_i$$

$$H \underline{w} = \sum_{i=1}^{i=N} \lambda_i (\underline{v}_i \underline{v}_i^T) \underline{w} = \sum_{i=1}^{i=N} \lambda_i P[\underline{v}_i] \underline{w} = \sum_{i=1}^{i=N} \lambda_i \underline{v}_i k_i$$



$$p(\lambda_i \underline{v}_i) = (k_i)^2$$

Probability of λ_i



Example (Continued):

$$\sum_{i=1}^{i=N} \lambda_i \underline{v}_i (\underline{v}_i^T \underline{w}) = \sum_{i=1}^{i=N} [\lambda_i (k_i)] \underline{v}_i$$

$$(k_i)^2: \left[\{|\underline{w}| = 1\} \Rightarrow \left\{ \sum_{j=1}^{j=N} (k_j)^2 = 1 \right\} \Rightarrow (k_i)^2 < 1 \right]$$

$$\left[\left\{ \sum_{j=1}^{j=N} (k_j)^2 = 1 \right\} \& (k_i)^2 < 1 \right] \Rightarrow$$

$(k_i)^2$ has the properties of a probability.

Example:

The energy of a particle can be **effectively** modelled by a Hermitian Operator.

The location and velocity states of the particle can be **effectively** modelled by sets of eigenvectors of the Hermitian operator.

The value of the particle's energy when in one of the states can be **effectively** modelled by the eigenvalue corresponding to that state.

The Heisenberg Uncertainty Principle says we can't simultaneously measure the location state/value and the velocity state/value from the 2 sets.

Vectors & QM

Quantum Mechanics is “just” modeling a physical system by “the right” Hermitian vector space.

Measurable Quantity $\xleftrightarrow{H^* = H}$ Hermitian Operator
Measured Value $\xleftrightarrow{H \underline{v} = \lambda \underline{v}}$ Eigenvalue

[All states = length 1, all eigenvalues real]

[Sometimes we don't care about the values!]

Measured State [“Pure”] $\xleftrightarrow{H(-\underline{v}) = \lambda(-\underline{v})}$ Eigenvector
[Sometimes we know the eigenvectors so we don't need the operator!]
[Most times we care only about the line, not the +- direction!]

Gen. System States \longleftrightarrow Eigenvector Combination

Probability of Value \longleftrightarrow (Length)² of projection on resultant eigenvector [≤ 1]

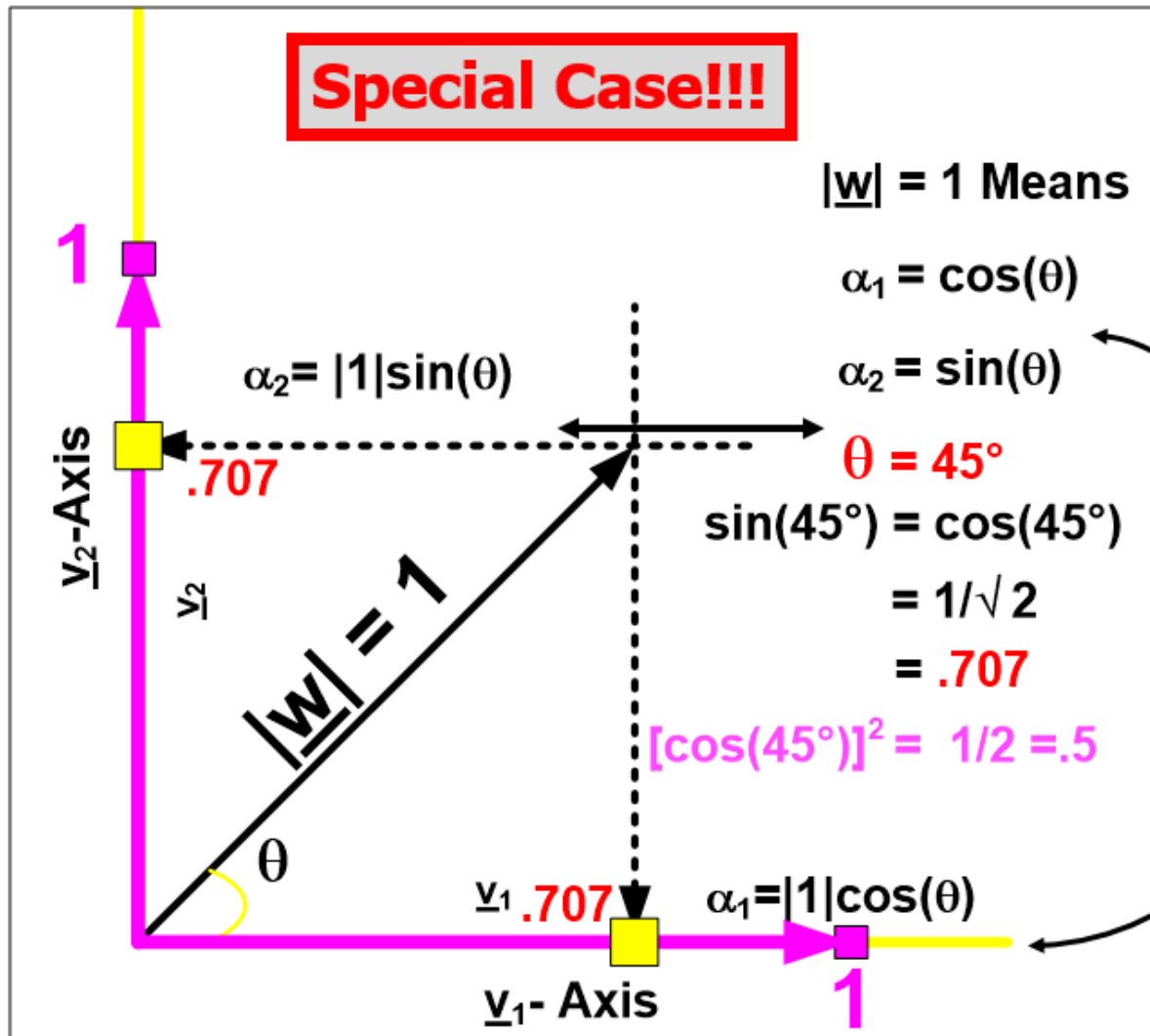
Vectors & QM

- 1. A Quantum System in a physical state is represented by a corresponding UNIT vector in some ~~abstract~~ Hermitian vector space.**
- 2. A Measurement puts the Quantum System into a unique physical state called a “Pure State” represented by a vector along a UNIT Basis Vector in that ~~abstract~~ vector space.**
- 3. Before any measurement, the system is in an unknown mixture of pure states, called a “Mixed State”.**
- 4. A measurement corresponds to a projection of a UNIT mixed vector onto ONE of the UNIT basis vectors of the ~~abstract~~ space.**

Vectors & QM

- 5. The $(\text{length})^2 < 1$ of the projected unit mixed state vector is the PROBABILITY of finding that Pure State in any given measurement.**
- 6. All Basis Vectors are actually eigenvectors of the operator representing the measured quantity.**
- 7. The value of the eigenvalue corresponding to the pure state is the measured VALUE in that pure state.**
- 8. A measurement corresponds to a meter reading of a physical system in an unknown state yielding a known state with a known metered value.**

Components as Projections



Basis Vectors

5. Physical Background of The QKD Algorithm

Background (2-D Polarization, & Probability)

1. 2-D Vector Uses

- Components as Projections

2. Polarization of Light

- Polarized Photons
- Filtered Photons Have $P=1$

3. Discrete Probability (Definition)

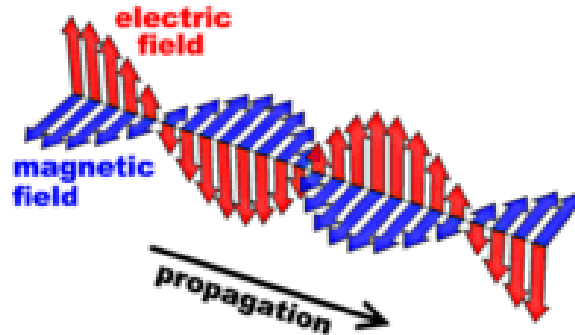
Addition ("OR")

Multiplication ("AND")

Physics of QKD

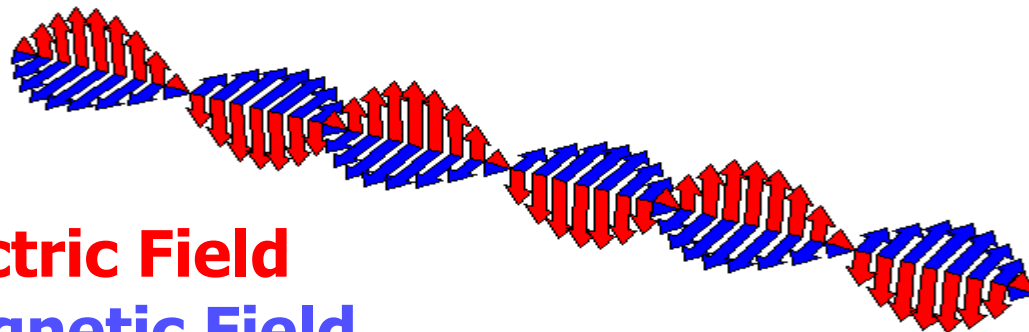
Propagating Electromagnetic Field

Rotates
Around
Direction
of
Propagation



IF Fixed Plane
THEN "Polarized"

2, 3-component Fields = 6 components



$\underline{E}(t)$ = Electric Field
 $\underline{H}(t)$ = Magnetic Field

Polarization of Light

- A photon is a “particle” of light
- A photon can be polarized along a direction

In the 2-D space perpendicular to the propagation direction.

- A photon can be polarized by a filter
- Once polarized by a filter (*QM Think*)



- it passes through that filter: $p = 100\%$
- it is blocked by a filter at 90° : $p = 0\%$

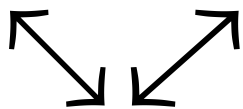
[2-D TRICK]

- it passes a (45°) filter BUT

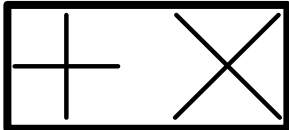
- it becomes (45°) polarized

- there is a 50% chance of being one

- there is a 50% chance of being other



2x(2-D Vector) Bases

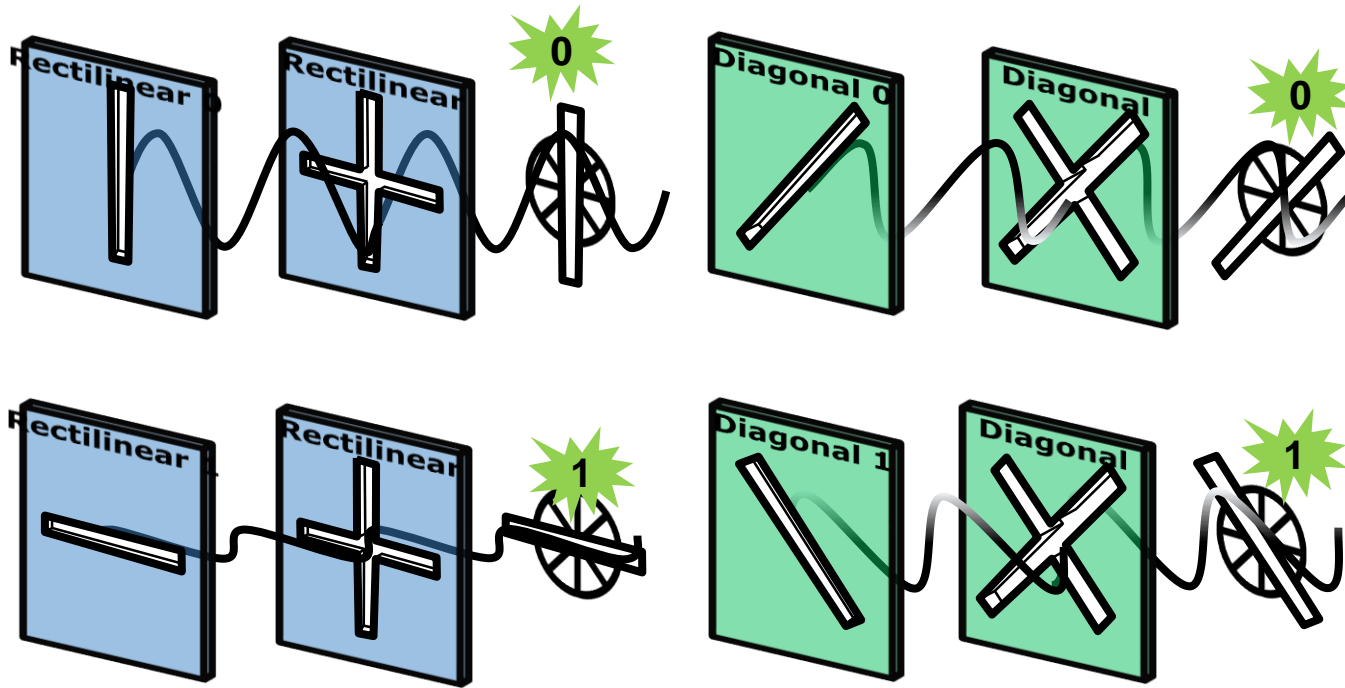
- A photon state is a unit 'vector' \updownarrow , \leftrightarrow , \swarrow , or \nwarrow .
[We use only the ray (line) not the direction]
- $\{ \updownarrow \ \& \leftrightarrow \}$ are a basis of the 2-D space
The 2-D space is perpendicular to the propagation direction.
- $\{ \swarrow \ \& \nwarrow \}$ are also a basis of the 2-D space
- These are also the 4 filters (*directions*) we use as
2 pairs $+$, \leftrightarrow & \times , \swarrow 
- A photon in a state in one basis is represented
 - as a sum in the other basis
 - with projected lengths = $1\cos(45^\circ) = 1/\sqrt{2}$
 - giving $[1\cos(45^\circ)]^2 = .5$ as probabilities

Polarized Photons

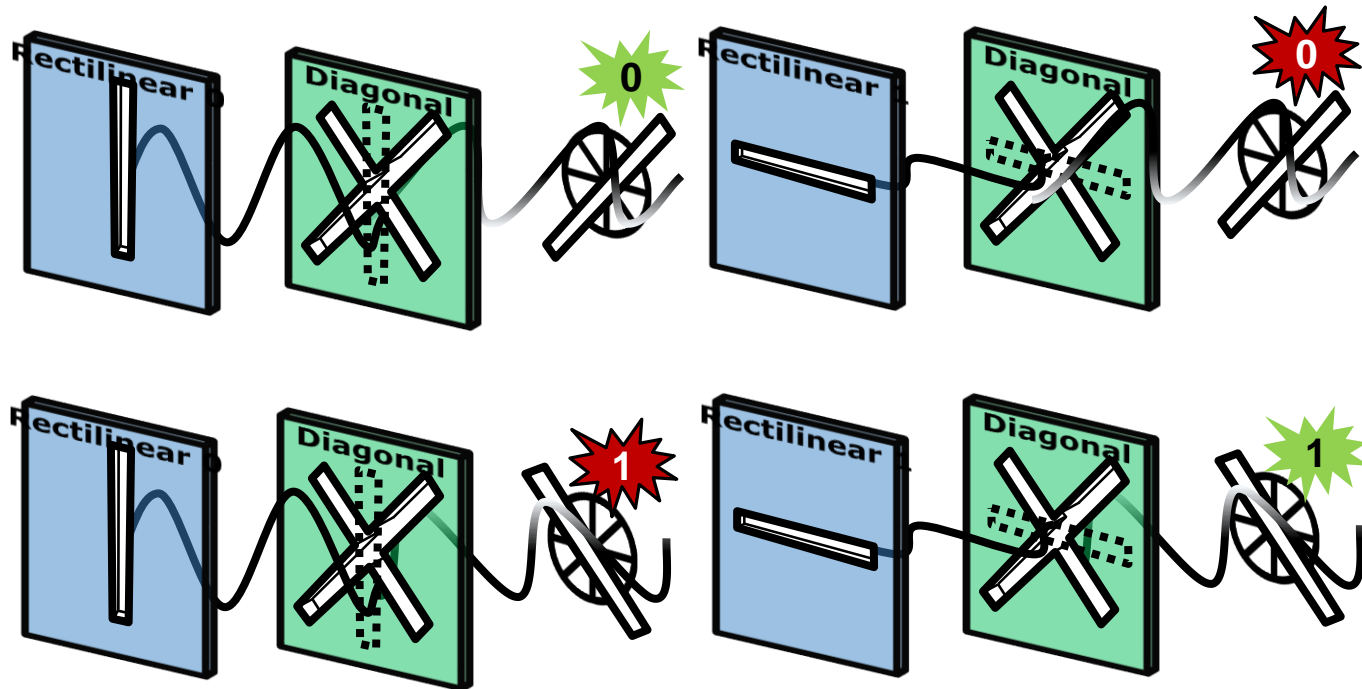
A UNIT basis vector represented in a Second, 45° rotated basis has (projection)² = .5 on EITHER second-basis direction.

I.e., we have $p=.5$ (I.e., EQUAL) probabilities of getting either second-basis vector as a measurement state result.

Polarization, correctly aligned filter correctly detects the bit sent

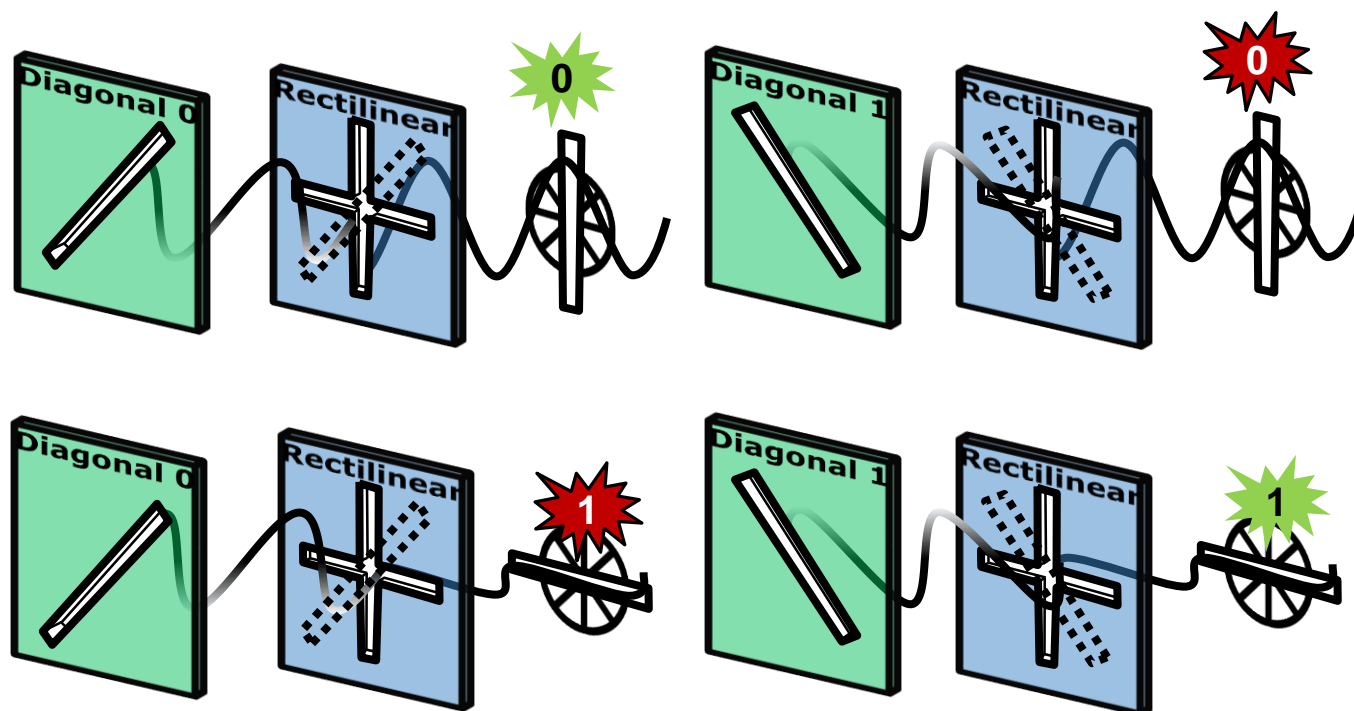


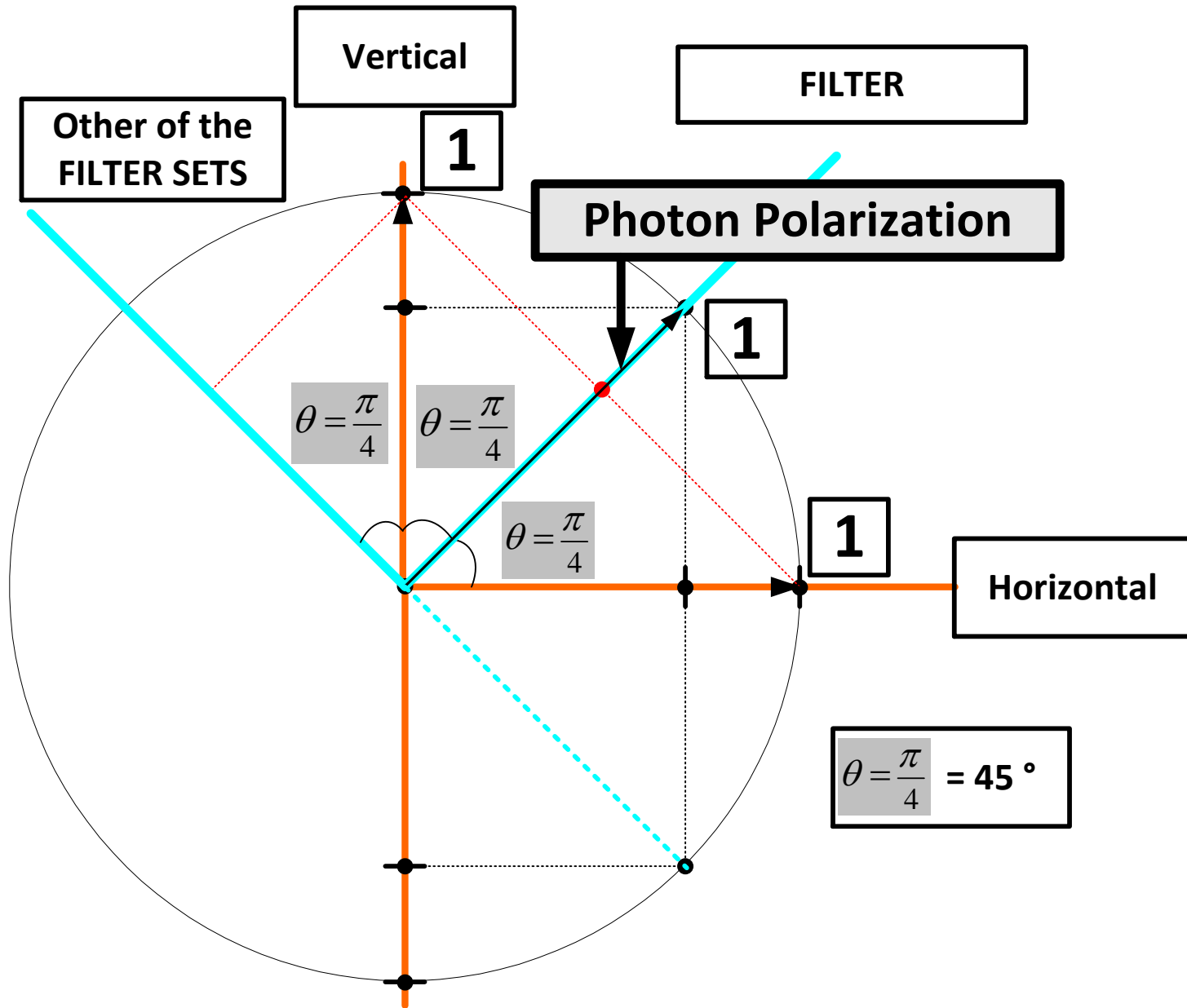
Rectilinear polarization, *diagonal* filter, quantum effect yields either bit



Photon has all possible states until detection, at which time it must choose a state based on the sending and detecting polarity filters

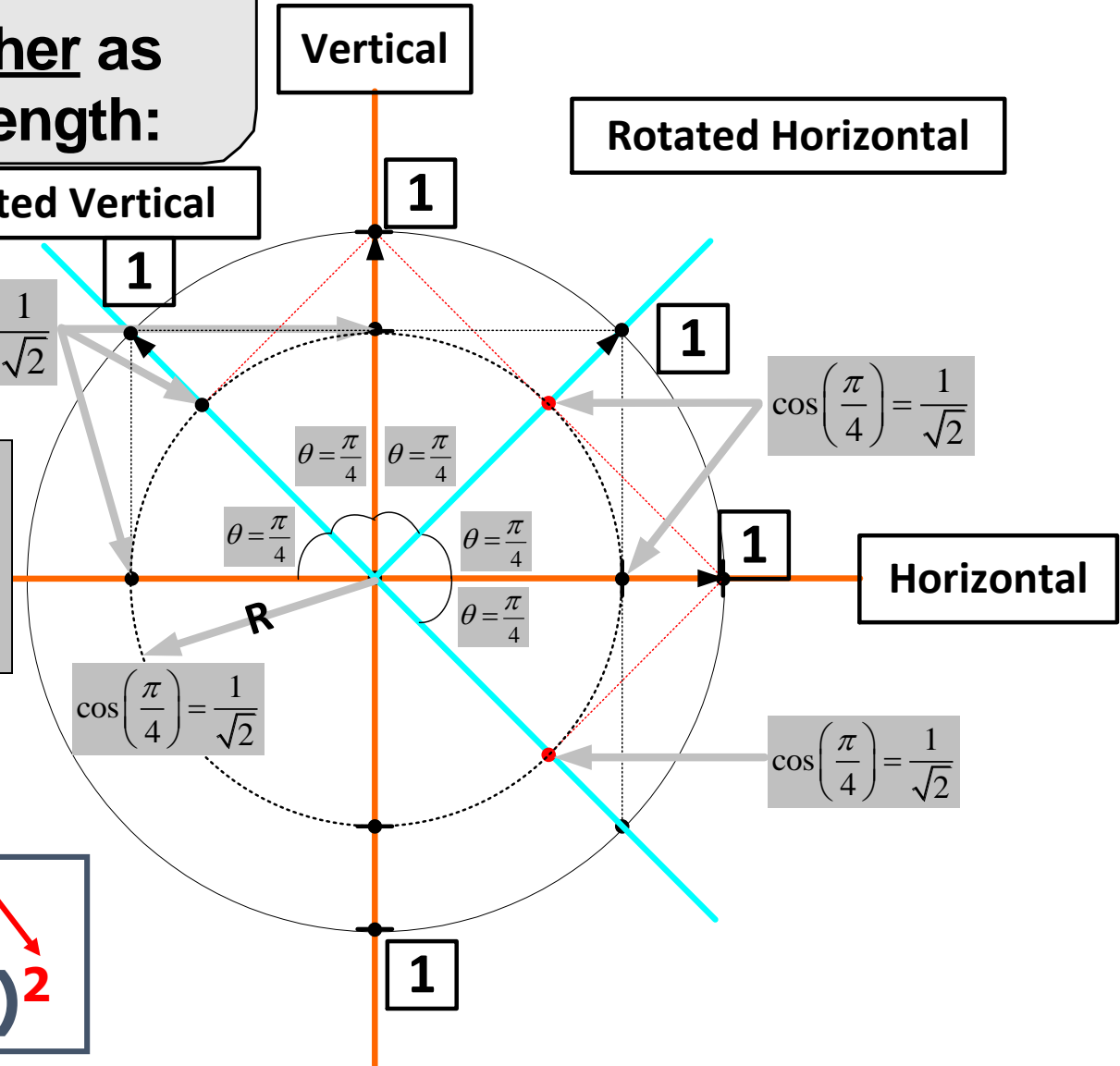
Diagonal polarization, *rectilinear* filter, quantum effect yields either bit





Two Axes Rotated 45 Degrees Relative to Each Other

The Unit Axis Vectors of Each Project onto the Other as Vectors of Equal Length:



Signs don't matter because we square lengths.

Rotated Basis
 $p = .5 = 1^2 \cos(45^\circ)^2$

**Filtered Photons Have $P=1$
Of passing same filter again.**

They lie along the filter direction

$$\text{so: } \cos(0^\circ)^2 = 1^2 = 1 = P.$$

OR

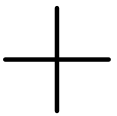
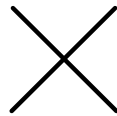
$$\cos(180^\circ)^2 = (-1)^2 = 1 = P.$$

So we care only about the line not the (+ or -) direction of the vector.

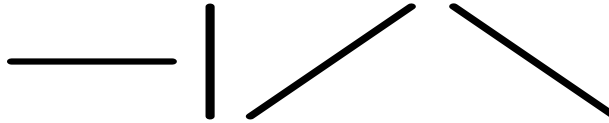
Discrete Probability

- $0 \leq p_i \leq 1$ for all cases i
- i discrete & finite
- {Sum of p_i over all cases i } = 1
- Probability of case j AND case k =
 $p_j p_k$
- Probability of case j OR case k =
 $p_j + p_k$

QKD Algorithm Background [1/3]

1. The whole purpose of the QKD algorithm is to find a secure 1TP (Key) for encryption
2. A polarized photon is a particle of light that has a known [i.e., measured] electric field orientation orthogonal to propagation
 - **PASSING LIGHT THRU A POLARIZATION FILTER IS MEASURING ITS FIELD ORIENTATION**
 - We use two filter **SETS** called   because that is what they look like.
 - They are rotated relative to each other by **45° TRICK!!! TRICK!!! TRICK!!!**

QKD Algorithm Background [2/3]



- We use 4 filters.
- We can call them horizontal, vertical, right 45 and left 45 (since the last two are at 45 degrees to the vertical).

- Polarized Photon State Vectors

- A state of vertical polarization is notated $|\updownarrow\rangle$
- A state of horizontal polarization is notated $|\leftrightarrow\rangle$
- A state of 45° right polarization is notated $|\nearrow\rangle$
- A state of 45° left polarization is notated $|\searrow\rangle$
[We really want only the ray not the direction since signs don't count because we square lengths]

QKD Algorithm Background [1/3]

We chose one state from each basis pair to represent a 1 bit
(the other of the pair is the 0 bit) **TRICK**

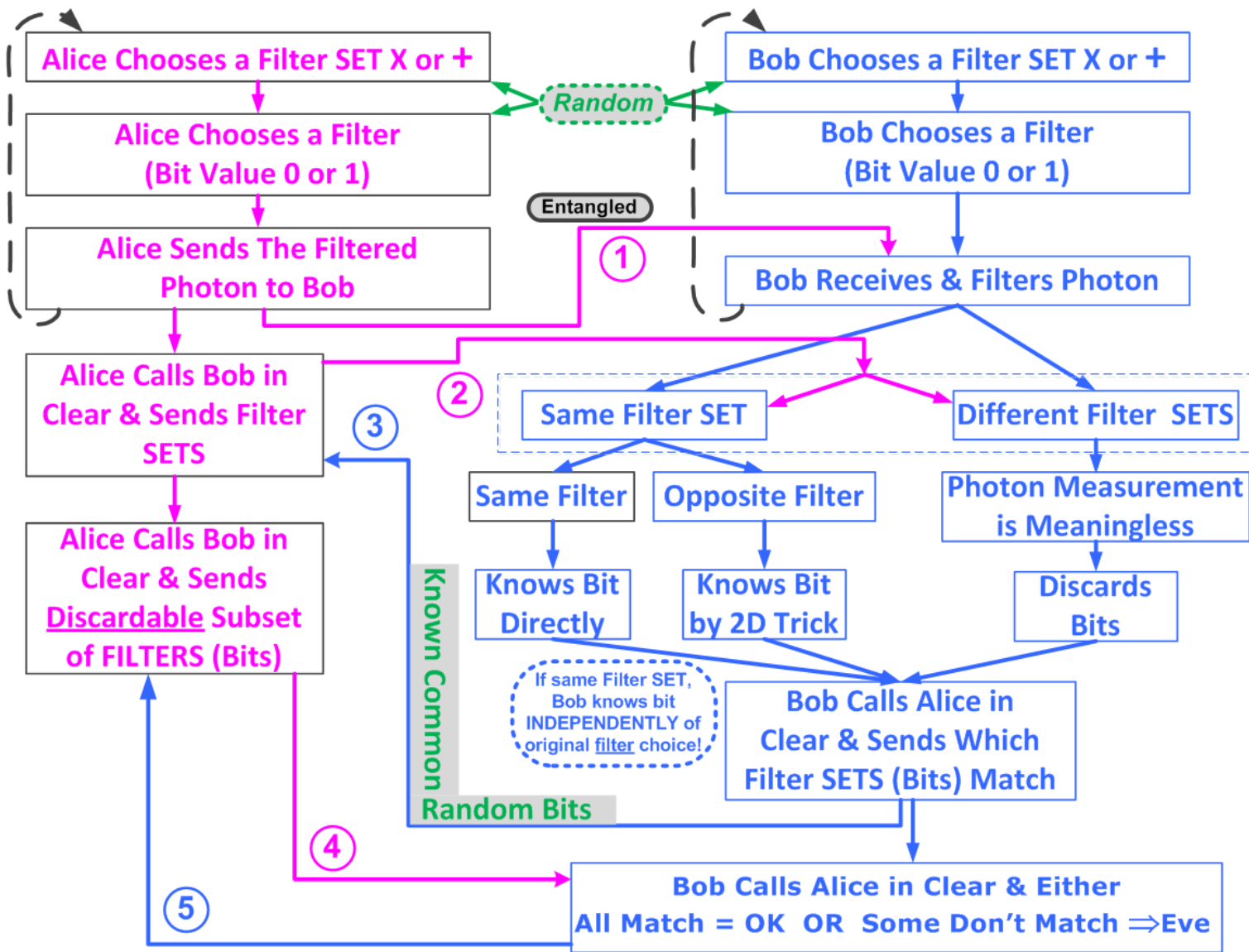
[Arbitrary choice of rep]

$$\begin{array}{l} |\updownarrow\rangle \triangleq 1, \quad \Rightarrow \quad |\leftrightarrow\rangle \triangleq 0 \\ |\nearrow\rangle \triangleq 1, \quad \Rightarrow \quad |\searrow\rangle \triangleq 0 \end{array}$$

6. Quantum Key Distribution The Algorithm

What QKD IS (The Details)

See the Bibliography for sources.



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QKD Algorithm

QKD_Alice-Bob_PhotonsV2.vsd

Quantum Key Distribution Algorithm

Alice sends N random bits (photons) using a random choice of N **filters**. Alice knows her bits (filters) & sets.

1. Bob uses a random choice of receiving **filters**
 - Bob knows his measured bits (filters) & sets.
 - Some are errors because he chose the wrong set
 - A bad set gives a bit error 50% of the time
 - A good set gives a correct bit 100% of the time
2. Alice calls Bob in the open and tells him **HER Filter SETS**
3. Bob tells Alice which of **HIS SETS agree (M bits)**
 - This determines a **secret set of M known bit values**
 - This is a key (after 4) for encryption - if no Eve
4. Alice calls Bob and reads to him a **discardable subset of HER actual FILTERS (i.e., BITS)**. If they agree there has been no Eve. Otherwise, there has been an Eve.
DISCARD ALL!

Entanglement
Goes Here

Singh, Simon. *The Code Book*. Anchor Books NY. ISBN 0-385-49532-3 (1999) PP. 339-344

Discussion [1/6]

Any measured (filtered) Photon is in a **pure** state

If measured again by the **same** filter get same state.

If measured by the other filter of the same **SET** (90°) see **NOTHING** so know bit (**2-D TRICK**).

If measured by the **OTHER SET**, get one of **them** with $p = .5$ by QM rep in other basis.

Notice neither Bob or EVE knows Alice's filters when they have to choose their own.

Discussion [2/6]

Bob's random choice of a **filter set matching Alice's** is equi-probable ($p = .5$).

Either **choice of bit (particular filter)**, given a matching pair will give correct info (actual bit or NOTHING, which implies the other **bit 2-D TRICK!!!**).

A choice of picking correct filter set ($1/2$).

The chance of picking N matching filters to Alice's hidden choices is $(1/2)^N$ [the AND case]. $(1/2)^N$ is $1/(2^N) \sim 10^{(-N/3.3)}$
For $N = 128 \sim 10^{(-36)}$

Doing it 100 times a second for one second $\sim 10^{(-36)100}$
 $\sim 10^{(-3600)}$

$\sim 10^{(-3600)}$ qualifies as the definition of impossible.

Discussion [3/6]

Individual Photon Polarization Measurement is a Quantum Process

Knowing that the wrong basis gives either result with $p = .5$ (**therefore no knowledge**) is a quantum result.

Knowing that $p = .5$ because of the probability law of mixed state projections in **45°** is a quantum result.

Knowing that a result of NO PHOTON means the complimentary pure state (therefore **full knowledge** - in **2-D ONLY**) is a quantum result. **2-D, IS A TRICK - AGAIN.**

Discussion [4/6]

The No Cloning Theorem

Any attempt to measure (read) an **unknown (mixed)** state **MUST** modify (Project) that state.

What we know is only the outcome state of the measurement, **not the input state.**

So – **we can't copy (clone) a state.**

The No Cloning Theorem

Discussion [5/6]

Quantum Key Distribution EVE Eavesdropping

Any attempt to read an unknown (mixed) photon and pass it on will introduce a probabilistic error.

There is a No Cloning Theorem.

In this case, cloning involves reading a photon.

Reading means applying a filter.

Eve can only pick a random choice of **filter & SET which introduces a random change to an incoming photon – sometimes - and sometimes not.**

She never knows which!

Discussion [6/6]

Quantum Key Distribution EVE Eavesdropping

Only if her filter **SET** happens to match the filter **SET** used by Alice to send the photon is there no error;

- Eve can't know if there is a match.
- A possible basis change causes ambiguity in her resultant measurement knowledge.
- No Cloning causes her to almost always pass on some changed photons.
 - **[She can be detected.]**

Algorithm Diagrams

- **UML Swim Lanes w/o Eve**
- **UML Swim Lanes w/ Eve**
- **Text Formulation**
- **Flowchart (Again)**

The Following Diagram: QKD Algorithm Overview

W/O Eve 1 of 2

W/O Eve 2 of 2

The The Following Diagram: QKD Algorithm Overview

W/ Eve 1 of 2

W/ Eve 2 of 2

Quantum Key Distribution Algorithm Summary

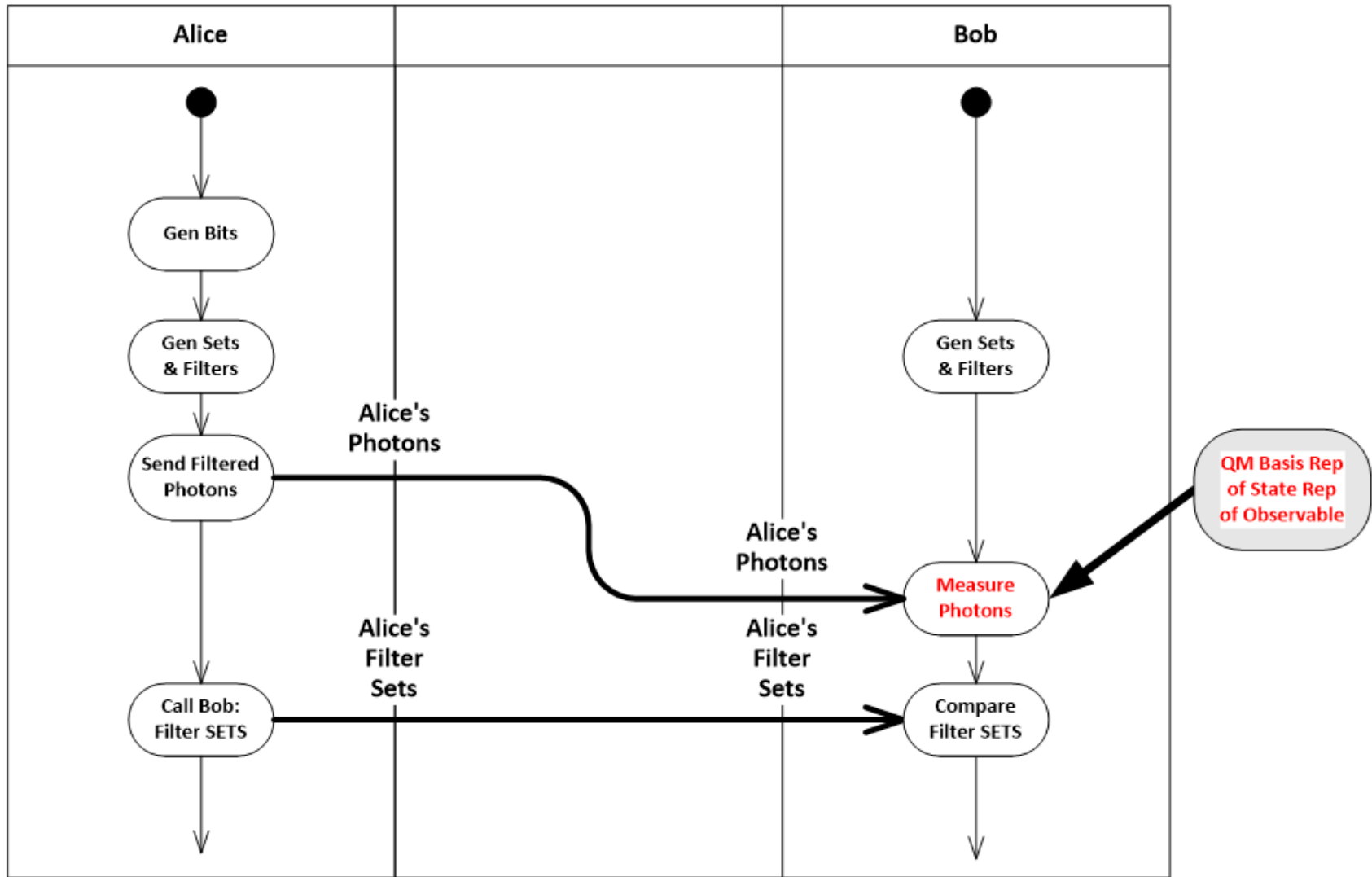
Quantum Key Distribution Results

Quantum Key Distribution One Time Pad

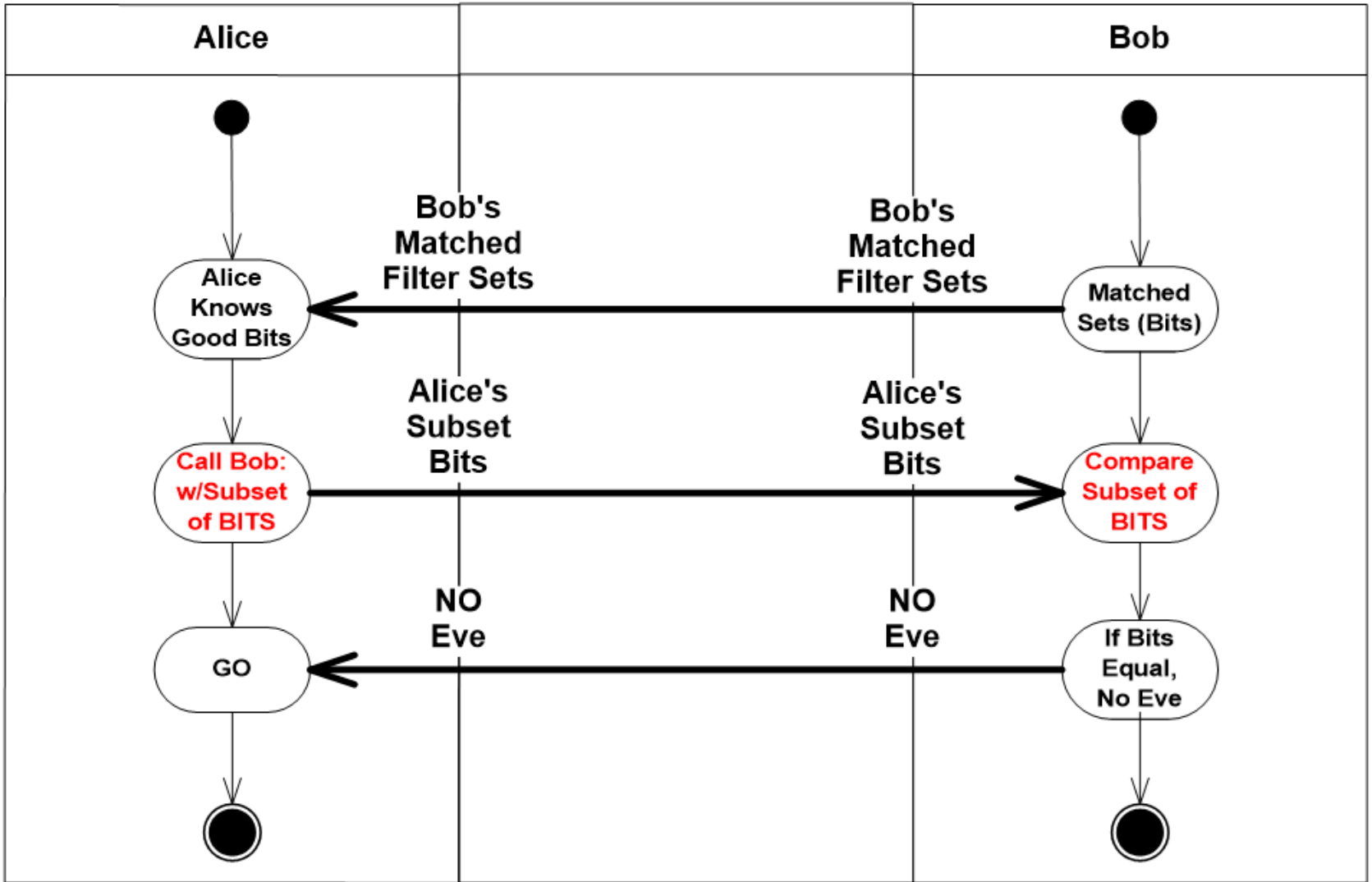
SUMMARY of Algorithm

- **We can securely transmit an unbreakable one time pad (Symmetric Key) of any desired length.**
- **We can ALWAYS detect EVE eavesdropping.**

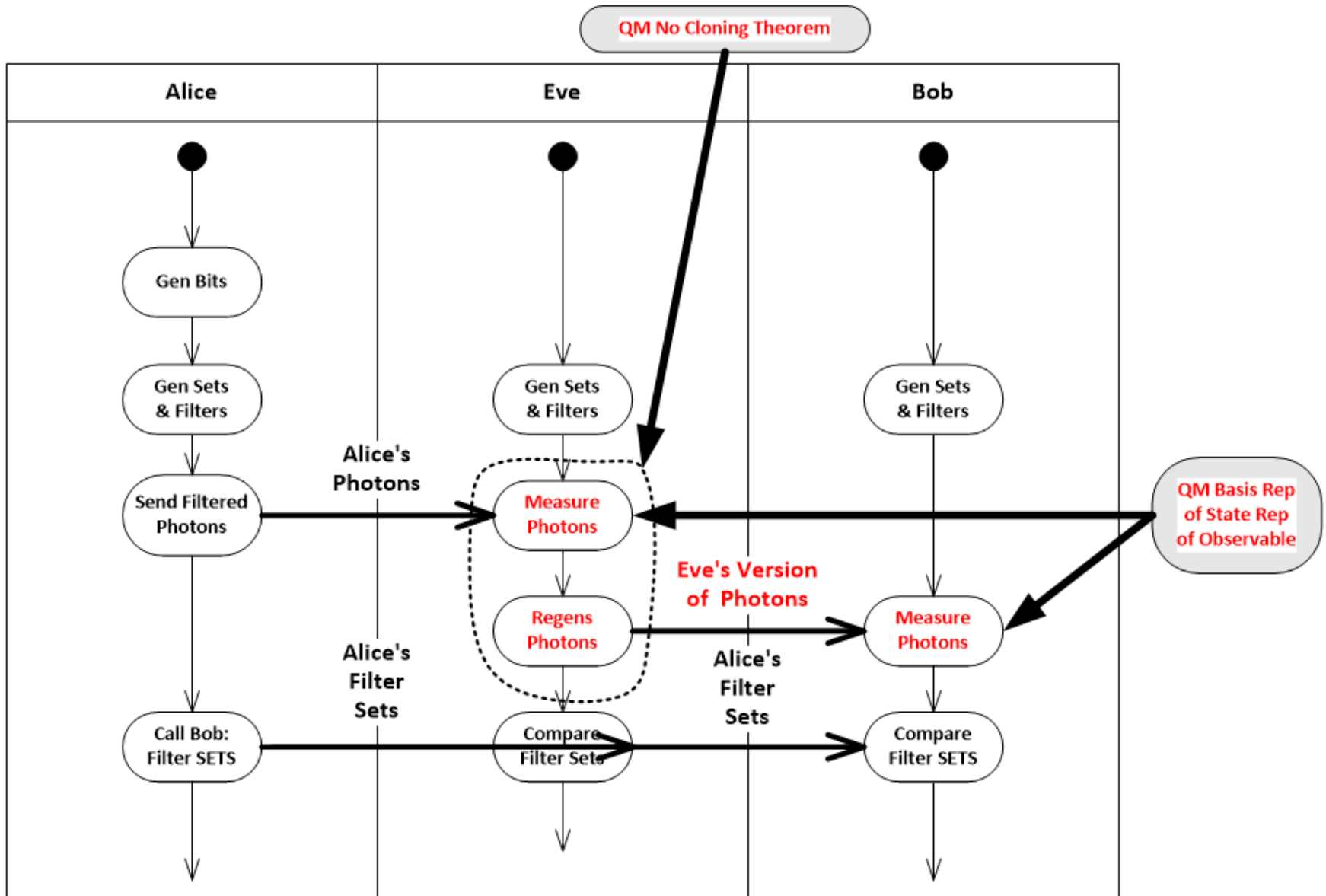
W/O Eve 1 of 2



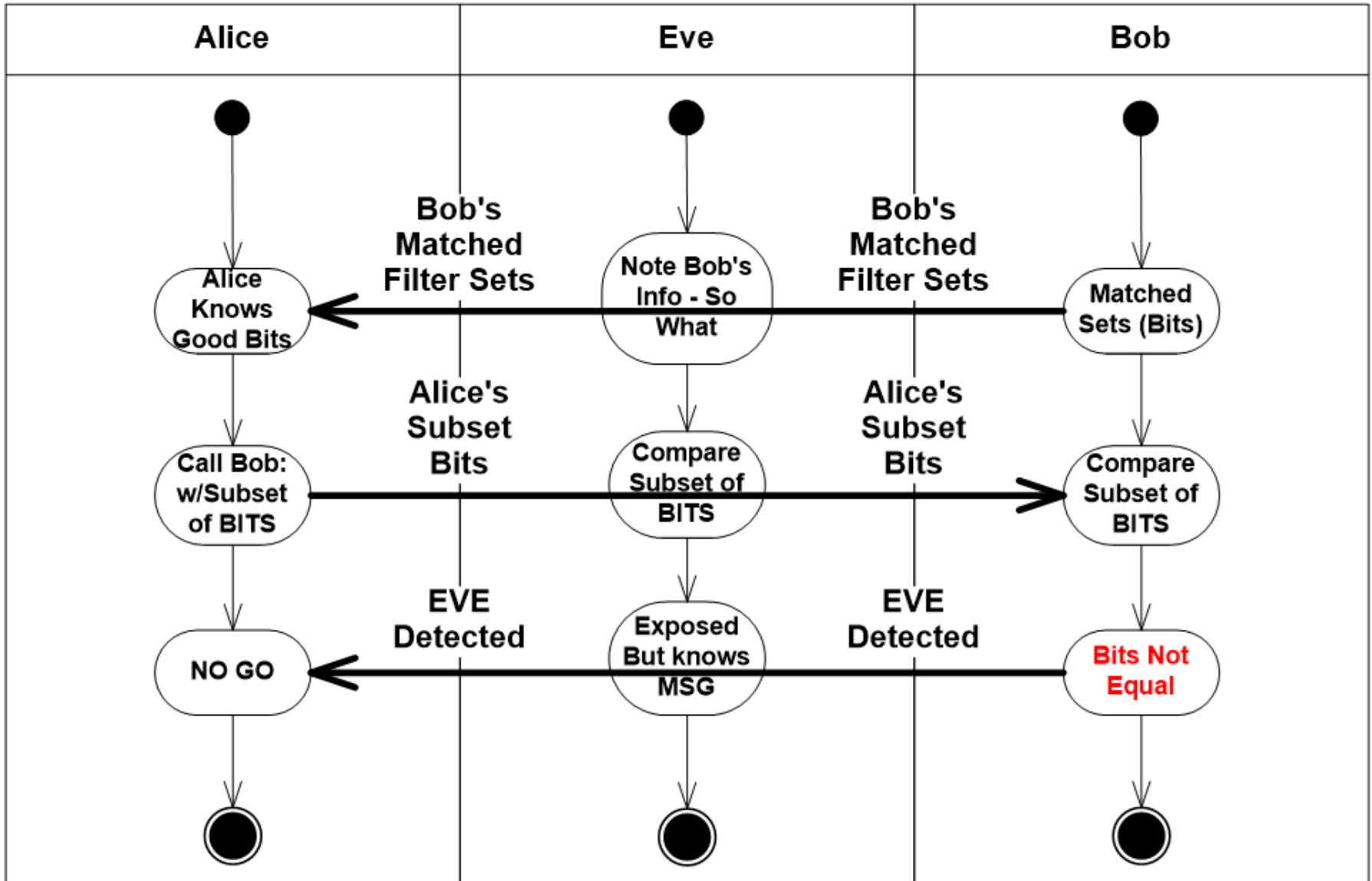
10. W/O Eve 2 of 2



W/Eve 1 of 2



W/Eve 2 of 2



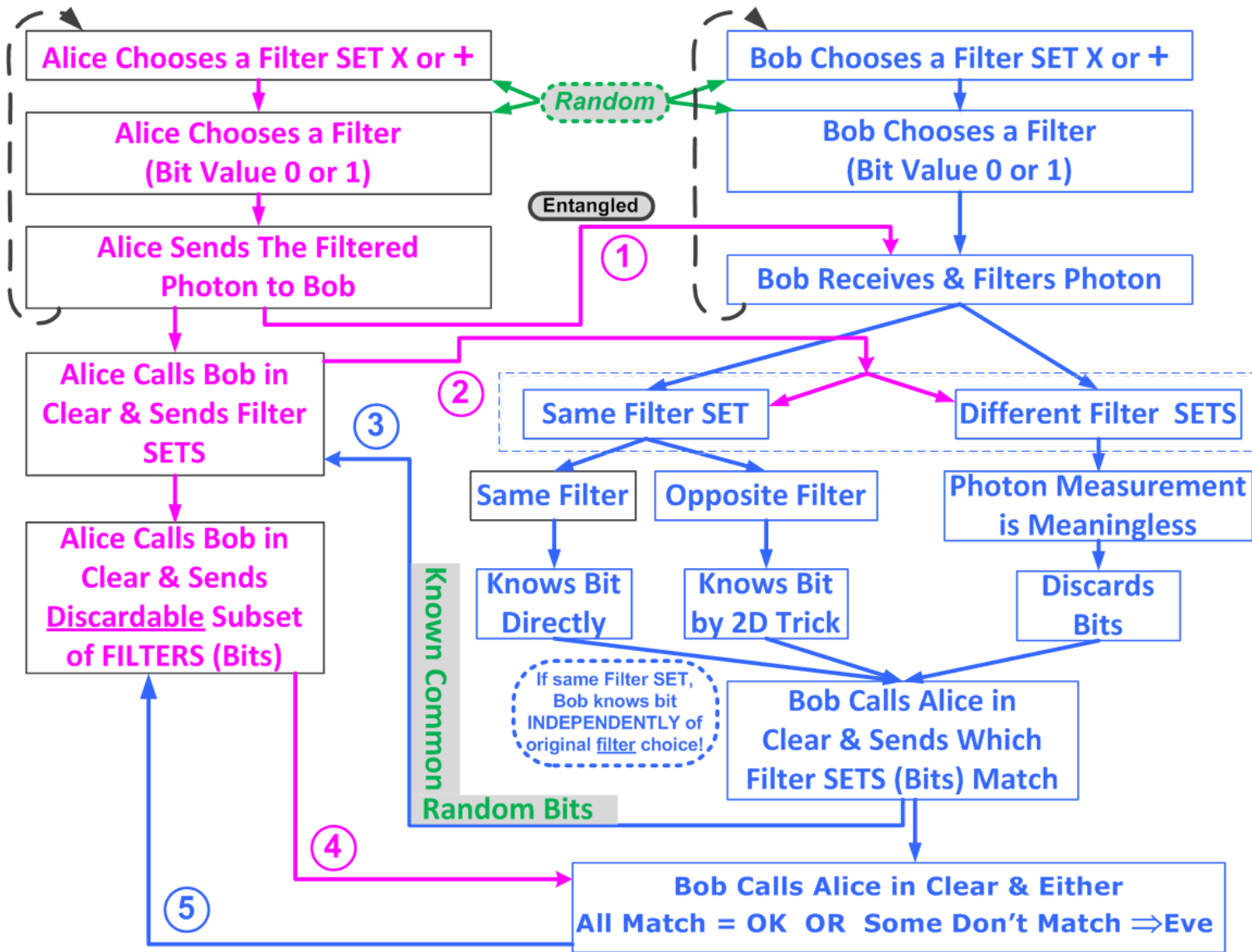
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Entanglement
Goes Here

Singh, Simon. *The Code Book*. Anchor Books NY. ISBN 0-385-49532-3 (1999) PP. 339-344



Quantum Key Distribution Results

- 1. This leaves both with a long random bit string which is secret and has not been read by an Eve ($P \sim 1$).**
- 2. This bit string is used as a secure symmetric key for a one-time-pad.**
- 3. The Navajo box generates new keys every 10 ms (100/sec).**

Quantum Key Distribution

1TP

**One-time-pads are
(classically)
known
(i.e., proven)
Unbreakable
(By Shannon).**

Algorithm Results

We can securely transmit an unbreakable one time pad (Symmetric Key) of any desired length.

We can ALWAYS detect EVE eavesdropping.

Algorithm Uses [1/2]

We can use the quantum key to **distribute** secure encrypted messages.

We can use the quantum key to **distribute** classical Private Keys (as messages).

BORMING's Dissertation

Algorithm Uses [2/2]

We can use the quantum key to distribute classical messages with a secure digital signature [Open text with encrypted hash of long message].

7. Quantum Key Distribution Bibliographies

Bibliographies

- 1. Quantum Economy**
- 2. Quantum Computing**
- 3. Polarization Demo Materials**
- 4. QKD**
 - 1/2**
 - 2/2**
- 5. IPSEC**
- 6. RIF**

Bibliography

(1. Quantum Economy)

**Waite, Stephen R., 2002,
Quantum Investing. Thomson Texere.
ISBN 1-587-99140-3**

(2. Quantum Computing)

Millburn, Gerard J., 1998, The Feynman Processor. Helix Books (Perseus Basic Books). ISBN 0-7382-0173-1.

{Layperson's introduction to Quantum Entanglement & Quantum Computing/Computers. Also introduces the difference between classical probability and quantum probability – which defines the quantum domain.}

Bibliography

(2. Quantum Computing)

Gudder, Stan “Quantum Computation.” **The American Mathematical Monthly** March 2003 # 110 pp. 181 – 201.

{Intro to QC for the mathematically prepared under grad.}

Johnson, George, 2003, A Short Cut Through Time. Alfred A. Knoph Pubs. ISBN 0-375-41193-3.

{Layperson’s introduction to QC & QE. My choice for selected topic readings in QC in IS courses.}

Bibliography

(3. Polarization Demo Materials)

B & H, <http://www.bhphotovideo.com/> 2004

{Search Binoculars & Scopes, 93608. (\$29.95). This is a Celestron Polarizing Lens Filter Set containing two rotating polarizing lenses in a threaded lens housing.}

Edmund Industrial Optics, <http://www.edmundoptics.com/> 2004

{Search KIT, then OPTICS DISCOVERY KIT (\$17.95). This is an American Optical Society of America classroom experiments kit – ages 10 – adult.}

Edmund Scientifics, <http://scientificsonline.com/> 2004

{Search 3038490, then POLARIZER EXPERIMENTERS KIT (\$19.95)}

Bibliography

(4. QKD) [1/2]

Nielsen, Michael A. and Isaac L. Chuang, 2000, Quantum Computation and Quantum Information. Cambridge University Press. ISBN 0-521-63503-9

{The definitive text on QE, QC and QI.}

{Possibly the most widely referenced textbook in QC, QI, and QE (encryption). It contains a review of QM for information people, the no cloning theorem, and the BB84 QKD protocol on which this presentation is based. }

Singh, Simon, 2002, The Code Book.

Anchor Books ISBN 0-385-49532-3

{Includes a layperson's chapter on modern QKD.

My choice for selected topic readings in QKD in IS courses.}

Bibliography (4. QKD) [2/2]

Tanenbaum, Andrew S. 2003, Computer Networks. Prentice Hall PTR. ISBN 0-13-066102-3
{Pp. 731 –734 Under One-Time Pads, Under Network Security}

Products for QKD

<http://www.idquantique.com/> 2004, 2010

<http://magiqtech.com/> 2004, 2010

<http://www.quantiki.org/> 2010

Quantum Entanglement

The Age of Entanglement: When Quantum Physics Was Reborn, by Louisa Gilder. Knopf; 1 edition (11- 11-2008) ISBN-10: 1400044170 ISBN-13: 978-1400044177
{Lay historical discussion of personalities, events, and ideas. A super read.}

Bibliography (5. IPSEC)

- **IP Security Protocol (ipsec) [Home page]**
<http://www.ietf.org/html.charters/ipsec-charter.html>
- **IPSEC Security Document Roadmap [1998]**
<http://www.ietf.org/rfc/rfc2411.txt>
- **Security Architecture for the Internet Protocol [1998]**
<http://www.ietf.org/rfc/rfc2401.txt>
- ***Internetworking with TCP/IP* by Douglas E. Comer. Vol. 1. 4th Ed. Prentice Hall (2000) ISBN 0-13-018380-6**
- **SSL is a Netscape de facto standard.**
<http://wp.netscape.com/eng/ssl3/>

Bibliography (6. RIF)

- Frank, R. I. (2003).

The Quantum Computing (QC), Quantum Encryption (QE), and Quantum Information (QI) Curriculum (Why? Now? Never?) *Information Systems Education Journal*, 1 (46). <http://isedj.org/1/46/>. ISSN: 1545-679X. (Also appears in [*The Proceedings of ISECON 2003: §2132*](#). ISSN: 1542-7382.)

[Argument for 3 topics inclusion in the IS curriculum]

- Frank, R. I. (2004).

An Outline of the Prerequisite Topics and Module for a Quantum Encryption (QE) Module in an IS Course. Proceedings Americas Conference on Information Systems (AMCIS August 2004) (Security track)
http://aisel.isworld.org/article_by_author.asp?Author_ID=5578 [Outline of an undergrad IS course component & prerequisites.]

Bibliography (6. RIF)

- Frank, R. I. (2005).

An IS Undergraduate Course Module on Quantum Key Distribution. *Information Systems Education Journal*, 3 (33). <http://isedj.org/3/33/>. ISSN: 1545-679X. (Also appears in [*The Proceedings of ISECON 2004: §2243*](#). ISSN: 1542-7382.)

[Detail of an undergrad IS course component on QKD only.]

8. Appendix on Vector Algebra And Hermitian Inner product Spaces

WHY Vectors (Linear Algebra)?

The world can be **effectively** modelled by Objects that have Observable States with Measureable Values with given probability.

States can be **effectively** modelled by vectors.

Objects can be **effectively** modelled by Operators on vectors.

Measureable values of the object can be **effectively** modelled by eigenvalues of the eigenvectors of the object.

WHY Vectors (Linear Algebra)?

The **probability** of finding the initial (before measurement) system in the final eigenstate i with measured eigenvalue i after measurement, can be **effectively** modelled as the square of the projection of the initial (before measurement) system vector onto the eigenvector i (found as the result of the measurement).

Vectors

Vectors

A set of thingies that ADD, and scalars (numbers) can multiply them. {**Vector: +, •** and **Scalar: +, -, *, /**}

Component Model $(x_1, y_1, z_1) \Rightarrow 3D$

$(x_1, y_1, z_1) + (x_2, y_2, z_2) = (x_3, y_3, z_3) \Rightarrow$
component-wise

addition

$s(x_1, y_1, z_1) = (sx_2, sy_2, sz_2) \Rightarrow$

component-wise scalar multiplication

Inner Product $(x_1, y_1, z_1) \bullet (x_2, y_2, z_2) =$ a **scalar**

vector/vector multiplication = $[x_1 x_2, + y_1 y_2, + z_1 z_2]$

ADD (component-wise component multiplication)

Length $\{(x_1, y_1, z_1) \bullet (x_1, y_1, z_1)\}^{.5} =$ a scalar ≥ 0

$\{x_1 x_2, + y_1 y_2, + z_1 z_2\}^{.5} \geq 0$

Vectors

- 1. A vector space is a collection of thingsies that add ($\underline{u} = \underline{v} + \underline{w}$), associate $\underline{u} + (\underline{w} + \underline{z}) = (\underline{u} + \underline{w}) + \underline{z}$, have an identity ($\underline{u} + \underline{0} = \underline{u}$), an additive inverse ($-\underline{u} + \underline{u} = \underline{0}$), and commute (+) ($\underline{w} + \underline{z} = \underline{z} + \underline{w}$).**
- 2. There is also a field of scalars that multiply them: $s\underline{u}$. This is scalar multiplication. Scalars in QM are complex #s.**
- 3. In addition to this scalar multiplication, there is a (vector • vector) multiplication called the scalar product (= the inner product = the dot product). It yields a scalar and is notated ($\underline{u} \cdot \underline{v}$) [“ \underline{u} dot \underline{v} ”] or $\langle u, y \rangle$.**

Vectors

- 1. A Basis of an (n-D) vector space is a set of n vectors which are:**
 - a. Linearly independent (Can't sum to one of them)**
 - &**
 - b. Span (generate) all vectors of the space**

- 2. We can always find an orthonormal basis which are:**
 - a. Of length 1**
 - b. Mutually orthogonal.**

Vectors

"Hermitian" Inner Product.

1. $\langle -, - \rangle$ [maps vectors to a scalar (COMPLEX NUMBER)]
2. $\langle \underline{u} + \underline{v}, \underline{w} \rangle = \langle \underline{u}, \underline{w} \rangle + \langle \underline{v}, \underline{w} \rangle$ [$\underline{u}, \underline{v}, \underline{w}$ are vectors]
3. $\langle \underline{u}, \underline{v} + \underline{w} \rangle = \langle \underline{u}, \underline{v} \rangle + \langle \underline{u}, \underline{w} \rangle$ **Bi-Linear**
4. $\langle s\underline{u}, \underline{v} \rangle = s \langle \underline{u}, \underline{v} \rangle$ [s is a COMPLEX NUMBER]
5. $\langle \underline{u}, s\underline{v} \rangle = s^* \langle \underline{u}, \underline{v} \rangle$ [$*$ is COMPLEX CONJUGATION]
6. $\langle \underline{u}, \underline{v} \rangle = \langle \underline{v}, \underline{u} \rangle^*$ [Conjugate Symmetric]
7. $\langle \underline{u}, \underline{u} \rangle = |\underline{u}|^2 \geq 0$ [= 0 iff $\underline{u} = \underline{0}$]

$$\langle H\underline{u}, \underline{v} \rangle = \langle \underline{u}, H^* \underline{v} \rangle = \langle \underline{u}, H\underline{v} \rangle \quad [\text{Hermitian Definition}]$$

[Eric W. Weisstein et al. "Hermitian Inner Product." From [MathWorld](http://mathworld.wolfram.com/HermitianInnerProduct.html)--A Wolfram Web Resource. <http://mathworld.wolfram.com/HermitianInnerProduct.html>]

If \underline{u} and \underline{v} are real: $\langle \underline{u}, \underline{v} \rangle = \langle \underline{v}, \underline{u} \rangle = \text{def } (\underline{u} \cdot \underline{v})$
 $\underline{u} \cdot \underline{v} = |\underline{u}| |\underline{v}| \cos(\text{angle between them}).$

Vectors

- $|\underline{u}| = (\underline{u} \cdot \underline{u}^*)^{.5}$ is a real number. "Length"
- **Hermitian** operators [$H=H^*$] (* is conjugate transpose) map vectors to vectors in the vector space. $H\underline{u} = \underline{v}$.
- **Eigenvectors** ("ownvectors") of an operator H are those vectors that H maps into multiples of themselves. $H\underline{u} = \lambda\underline{u}$. If $|\underline{u}| = 1$, λ is an **eigenvalue** of H associated with \underline{u} [there can be more than one \underline{u} for a given λ].
- *An Hermitian operator's eigenvalues are real.*
- An Hermitian operator's eigenvectors form a **basis** of the entire [Hermitian Vector] space.
- In a real inner product space the symmetric operators ($A=A^t$) are the Hermitian operators.

9. Appendix on Sin & Cos

**Sin(nx) & Cos (nx)
Form an Orthonormal Basis
of an
Infinite Dimensional Space
(all n) and are the
Eigenvectors of the
Second Derivative Operator**

$$\frac{d^2(\quad)}{dx^2}$$

Mixed (n, m) sin(nx) / cos(mx) are orthogonal
Same (n = m) sin(nx) / cos(nx) are orthogonal

$$\int_{-\pi}^{\pi} \cos(nx) \sin(mx) dx = 0$$

Mixed (n, m) Sin(nx) / sin(mx) are orthogonal
Same (n = m) Sin(nx) / sin(nx) are normalizable

$$\int_{-\pi}^{\pi} \sin(nx) \sin(mx) dx = \pi \delta_{mn} \quad [m, n \geq 1]$$

i.e. $\int_{-\pi}^{\pi} \sin^2(nx) dx = \pi \quad [n \geq 1]$

i.e. $\int_{-\pi}^{\pi} \sin(nx) \sin(mx) dx = 0 \quad [m \neq n]$

Mixed $\cos(nx)$ / $\cos(mx)$ are orthogonal
Same $\cos(nx)$ / $\cos(nx)$ are normalizable

$$\int_{-\pi}^{\pi} \cos(nx) \cos(mx) dx = \pi \delta_{mn} \quad [m, n \geq 1]$$

i.e. $\int_{-\pi}^{\pi} \cos^2(nx) dx = \pi \quad [n \geq 1]$

i.e. $\int_{-\pi}^{\pi} \cos(nx) \cos(mx) dx = 0 \quad [m \neq n \ \& \ \geq 1]$

**cos(nx) & sin(nx) are
Orthogonal to a constant (1).**

$$\int_{-\pi}^{\pi} (1) \sin(mx) dx = 0$$

$$\int_{-\pi}^{\pi} (1) \cos(nx) dx = 0$$

So What?

Any

- **Continuous function $f(x)$ on $[-\pi, \pi]$**
- **With only a finite number**
- **of**
- **Finite jump discontinuities**

Equals the infinite sum

$$f(x) = \frac{a_0}{2} + \sum_{n=1}^{\infty} a_n \cos(nx) + \sum_{n=1}^{\infty} b_n \sin(nx)$$

where

$$\begin{aligned} a_0 &= \left(\frac{1}{\pi}\right) \int_{-\pi}^{\pi} f(x) dx = 0 & P_1[f(x)] &= \langle 1, f(x) \rangle \\ a_n &= \left(\frac{1}{\pi}\right) \int_{-\pi}^{\pi} f(x) \cos(nx) dx = 0 & P_{\cos(nx)}[f(x)] &= \langle \cos(nx), f(x) \rangle \\ b_n &= \left(\frac{1}{\pi}\right) \int_{-\pi}^{\pi} f(x) \sin(nx) dx = 0 & P_{\sin(nx)}[f(x)] &= \langle \sin(nx), f(x) \rangle \end{aligned}$$

$$\begin{aligned}\sin(nx) = -\sin(-x) &\Rightarrow \frac{d(\sin(nx))}{dx} = n \cos(nx) \Rightarrow \frac{d^2(\sin(nx))}{dx^2} = -n^2 \sin(nx) \\ \cos(nx) = \cos(-nx) &\Rightarrow \frac{d(\cos(nx))}{dx} = -n \sin(nx) \Rightarrow \frac{d^2(\cos(nx))}{dx^2} = -n^2 \cos(nx)\end{aligned}$$

$$H(\bullet) \approx \frac{d^2(\bullet)}{dx^2} \Leftrightarrow H \begin{bmatrix} \sin(nx) \\ \cos(nx) \end{bmatrix} = \begin{bmatrix} -n^2 \sin(nx) \\ -n^2 \cos(nx) \end{bmatrix}$$

**Therefore sin & cos are the Eigenvectors
of the
Second Derivative Operator**

